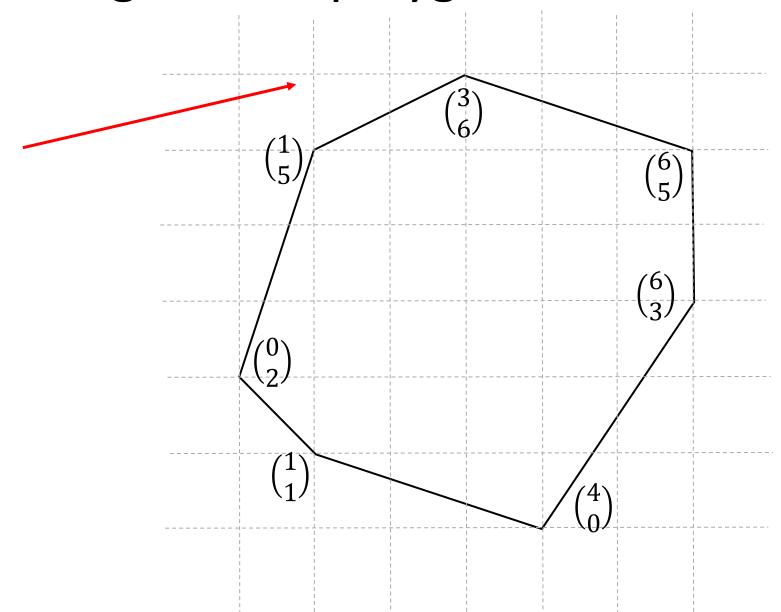
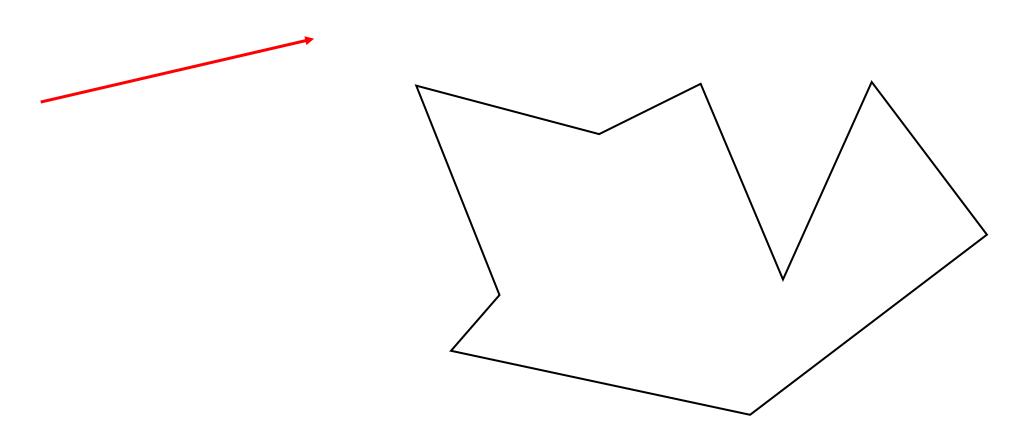
## Bisecting convex polygons



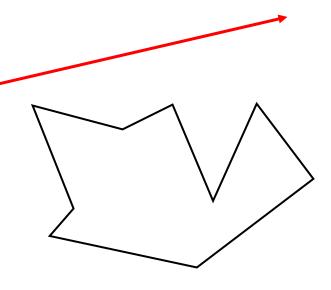
#### The problem I would really like to solve:

Given any simple connected polygon P (specified in terms of its coordinates) and a direction vector  $\mathbf{u}$ , write down the equation of the straight line bisecting the area of P in the direction of  $\mathbf{u}$ .

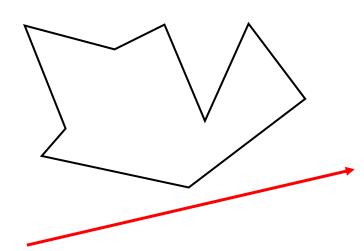


#### Which we know can be done...

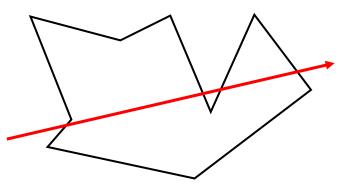
... by the intermediate value theorem...



Area 'above' line = 0



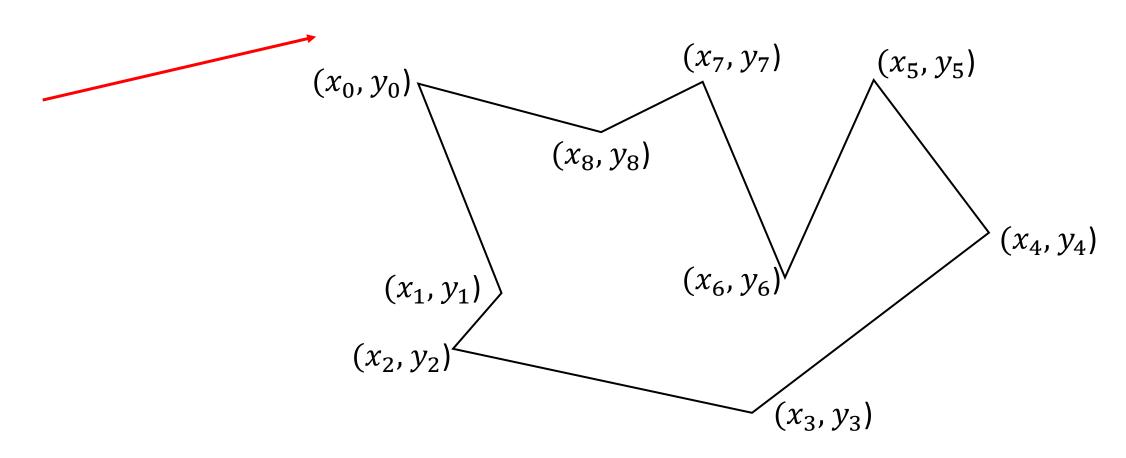
Area 'above' line = A (total area)



Area 'above' line = A/2

#### A restricted problem

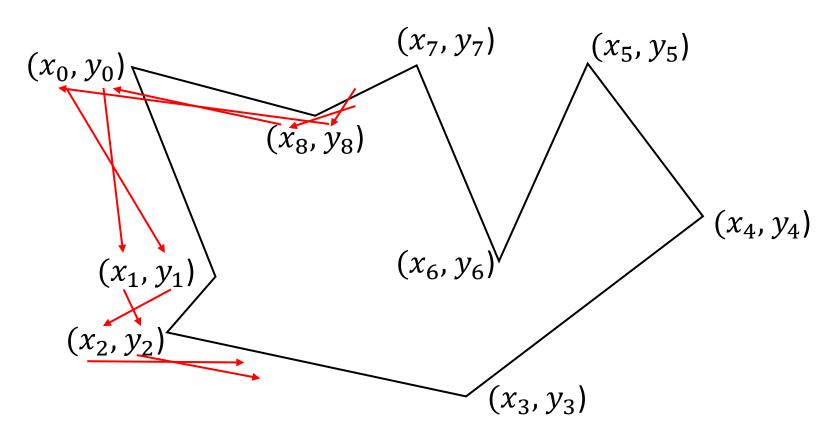
Given any simple connected polygon P, specified in terms of **rational** coordinates  $(x_i, y_i)$ , and a direction vector  $\mathbf{u}$ , specified as a **rational number** (its slope), write down the equation of the straight line bisecting the area of P in the direction of  $\mathbf{u}$ .



#### The shoelace formula

Given any simple connected polygon P, specified in terms of coordinates  $(x_i, y_i)$ , the area A of P is given by:

$$\frac{1}{2}(x_0y_1 - x_1y_0 + \dots + x_{n-1}y_0 - x_0y_{n-1}).$$

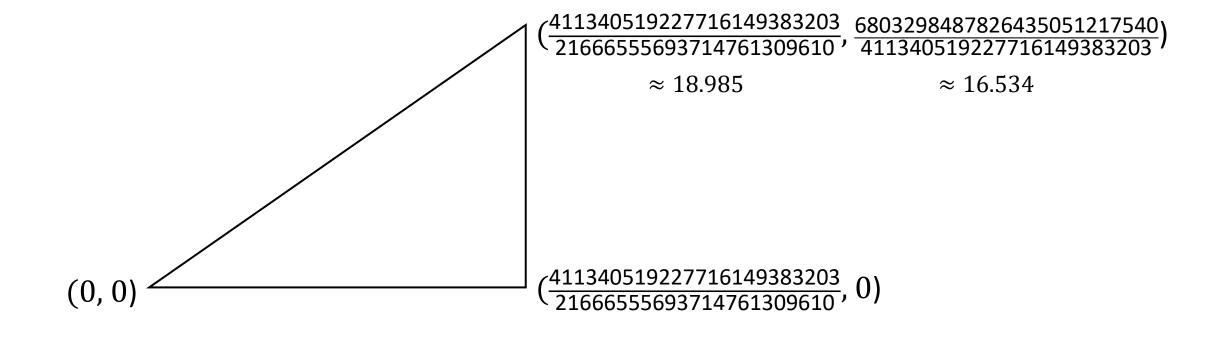


#### IVT and constructiveness

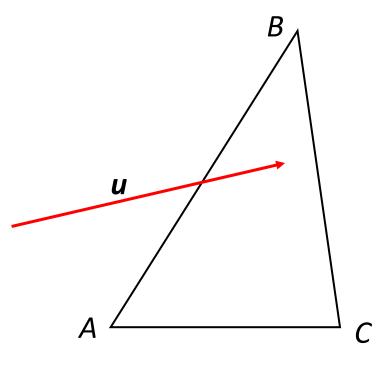
The shoelace formula says our restricted problem is solved by a straight line equation with rational coefficients.

But IVT will not necessarily produce this equation, only a rational solution arbitrarily close to it.

Area problems in the rationals are not necessarily straightforward! E.g. (Don Zagier) the 'simplest' right triangle with area 157 is:



#### Triangles: the easy case

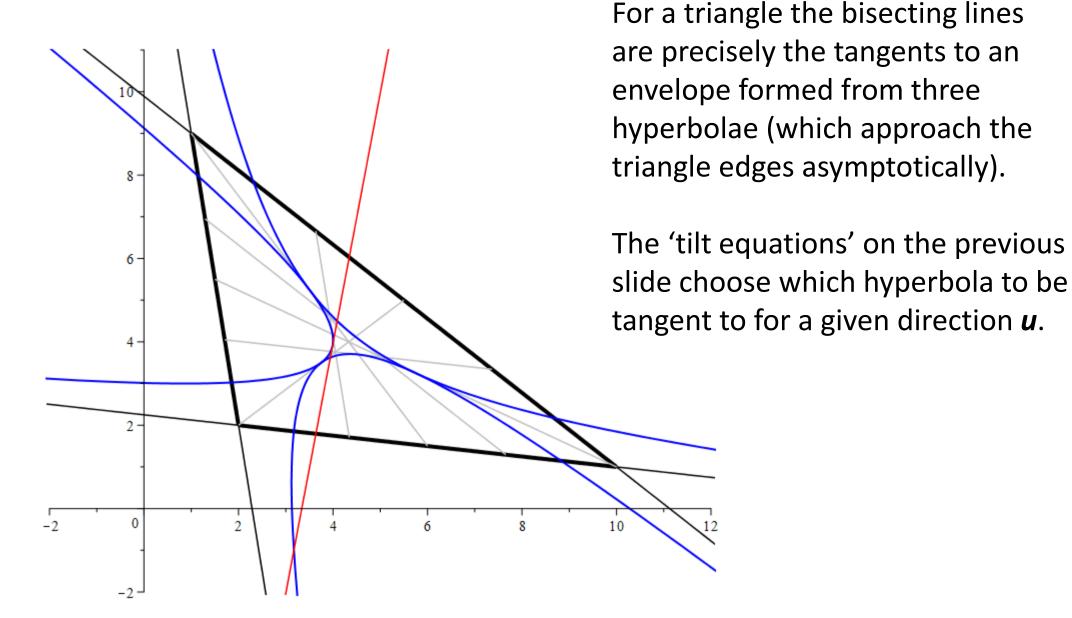


Bisecting straight line:

$$r = u\lambda + \begin{cases} (1-m)A + (1-t)mC + tmB & \text{for tilt on } BC \\ (1-m)B + (1-t)mA + tmC & \text{for tilt on } CA, \\ (1-m)C + (1-t)mB + tmA & \text{for tilt on } AB \end{cases}$$
 where  $t = \left(1 + \sqrt{\frac{2-w}{1+w}}\right)^{-1}$ ,  $m = 1/\sqrt{2}$ ,

and w is unique solution in [0,1] to tilt equations.

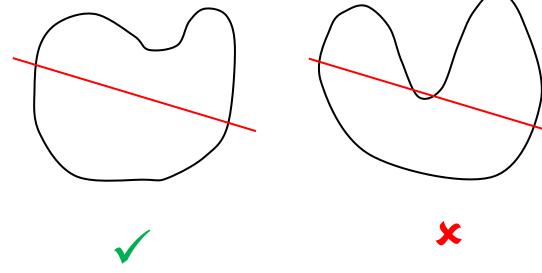
#### Bisection envelopes (triangles)



#### Bisection envelopes (polygons)



**Bisection-convex:** any bisecting straight line intersects the curve in exactly two points



Bisection envelopes

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**Proposition 3.3.** The bisection envelope  $\mathcal{B}$  of a polygon  $\mathcal{S}$  is the union of a finite number of sections of hyperbolas. Let the set of all asymptotes of these hyperbolas be H, and let the set of all lines that contain the sides of  $\mathcal{S}$  be G. Then  $H \subseteq G$ , with equality if no two lines in G are parallel.

#### Strictly bisection-convex curves

We now restrict the class of curves S to be studied.

**Definition 2.2.** Define S and L as above. We say that S is *bisection convex* if for all  $\theta$ ,  $l_{\theta}$  intersects S in exactly two points. Alternatively, for every point A on S, there exists a unique point B also on S such that the line AB bisects the interior area of S.

We also create a tighter restriction.

**Definition 2.3.** Define S and L as before. We say that S is *strictly bisection convex* if it is bisection convex and for all  $\theta$ ,  $l_{\theta}$  is not tangent to S. At any point where there are two tangents to S—one from each side—the  $l_{\theta}$  through that point is distinct from both tangents.

Henceforth, unless otherwise stated, it is assumed that S is strictly bisection convex.

Bisection envelopes

#### That IVT 2-pancakes issue again...

Define  $A(\theta)$  and  $B(\theta)$  to be the endpoints of the bisecting chord in direction  $\theta$ , with  $B(\theta) = A(\theta + \pi)$ . We distinguish between  $A(\theta)$  and  $B(\theta)$  by demanding that for each point  $Q \neq A(\theta)$ ,  $B(\theta)$  on the bisecting chord, the vector  $A(\theta) - Q$  points in positive direction  $\theta$  and the vector  $B(\theta) - Q$  points in positive direction  $\theta + \pi$ .

**Proposition 2.4.** Assume that S is bisection convex. Then  $A(\theta)$  varies continuously with  $\theta$ .

*Proof.* First, we note that any two bisecting chords must intersect in the interior of S, for if they did not, the interior of S would be split into three regions, one of which would have zero area, which does not make sense.

From this, we have  $\lim_{\epsilon \to 0} l_{\theta+\epsilon} = l_{\theta}$ , as the limit of the intersection point  $l_{\theta+\epsilon} \cap l_{\theta}$  is bounded. This also implies that the limit as  $\epsilon \to 0$  of the distance from  $A(\theta+\epsilon)$  to the intersection point  $l_{\theta+\epsilon} \cap l_{\theta}$  is bounded. Therefore, the limit as  $\epsilon \to 0$  of the perpendicular distance from  $A(\theta+\epsilon)$  to  $l_{\theta}$  is zero.

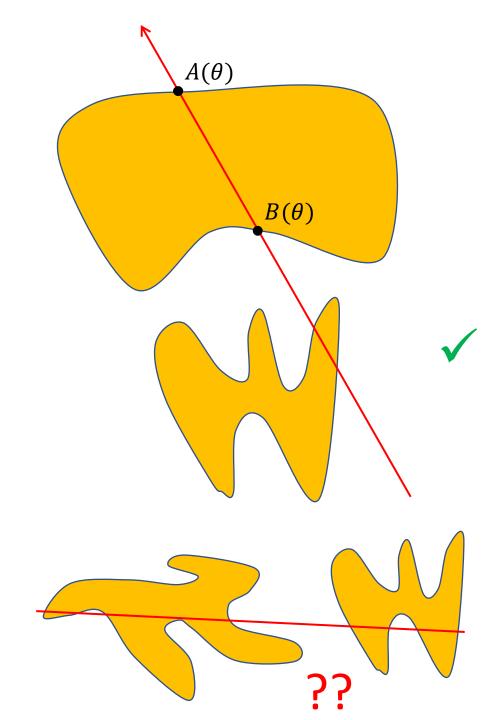
We have that  $\lim_{\epsilon \to 0} A(\theta + \epsilon)$  must be a point P on  $l_{\theta}$  which intersects S, where for every other point Q on the bisecting chord with direction  $\theta$ , the vector P - Q points in positive direction  $\theta$ . There is only one such point,  $A(\theta)$ ; therefore,

$$\lim_{\epsilon \to 0} A(\theta + \epsilon) = A(\theta),$$

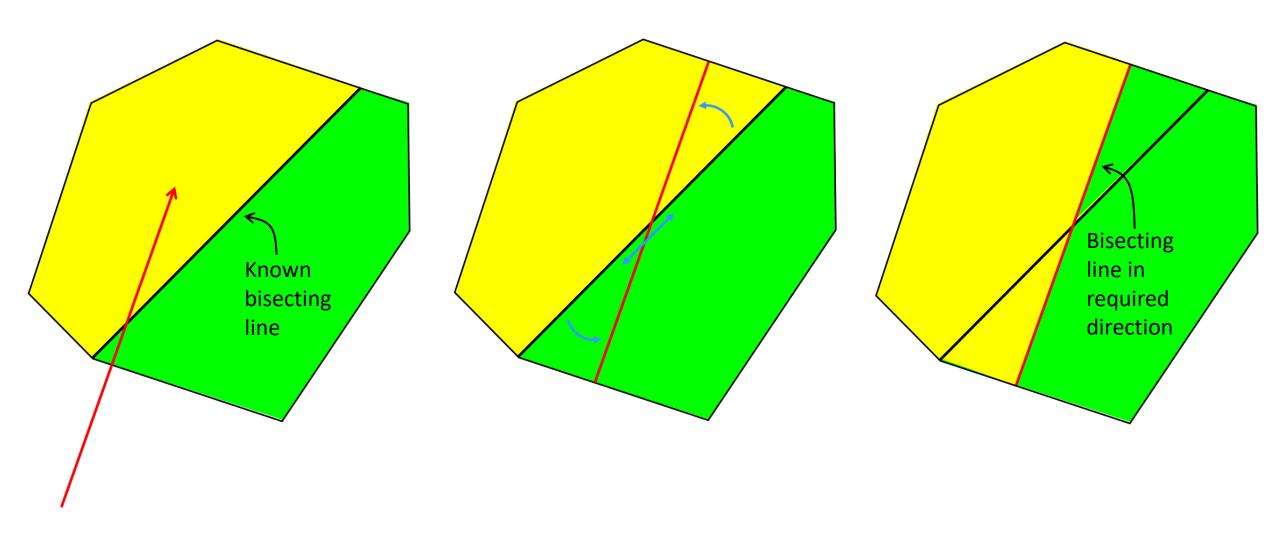
and  $A(\theta)$  varies continuously with  $\theta$ .

Bisection envelopes
Noah Fechtor-Pradines

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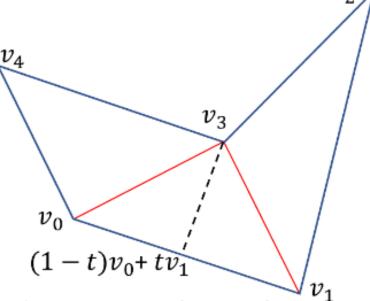
### A straight line equation for convex polygon bisection...



#### ... or even bisection-convex polygon bisection?

#### An Application

We may triangulate a polygon on n vertices by adding n-3 diagonals, as illustrated on the right. We would like to test if some straight line joining a triangle vertex to the opposite polygon edge bisects the area of the polygon. In our diagram this requires a value of  $t \in [0, 1]$  for which the poly-



gons  $v_0$ ,  $(1-t)v_0 + tv_1$ ,  $v_3$ ,  $v_4$  and  $(1-t)v_0 + tv_1$ ,  $v_1$ ,  $v_2$ ,  $v_3$  have equal area.

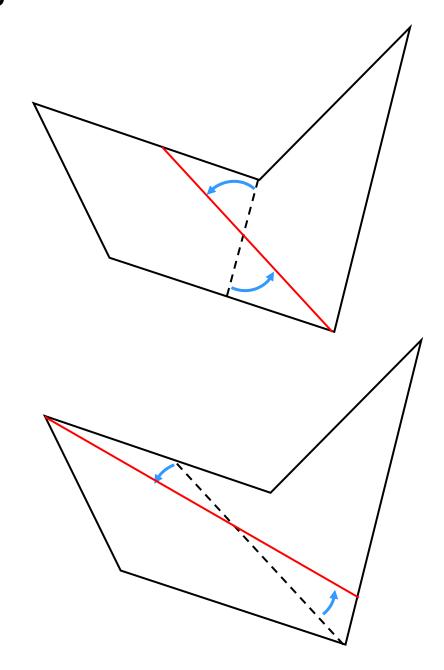
An application of the shoelace formula gives

$$t = \frac{A_R - A_L}{2A_{\Lambda}}$$

 $A_L$  = area to left of middle triangle

 $A_R$  = remaining polygon area

 $A_{\Delta}$  = area of middle triangle



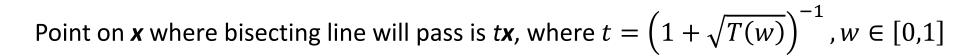
# Strategy for writing down a bisecting straight line equation

Divide interior of polygon into sectors bordered by pairs of bisecting lines x and y



$$u(x, y) = wx + (1 - w)y, w \in [0,1]$$

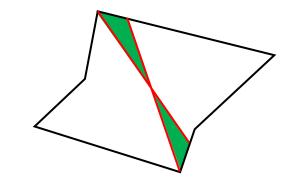
If given direction vector  $\mathbf{u}$  lies in sector then this will solve to give  $w \in [0,1]$ 



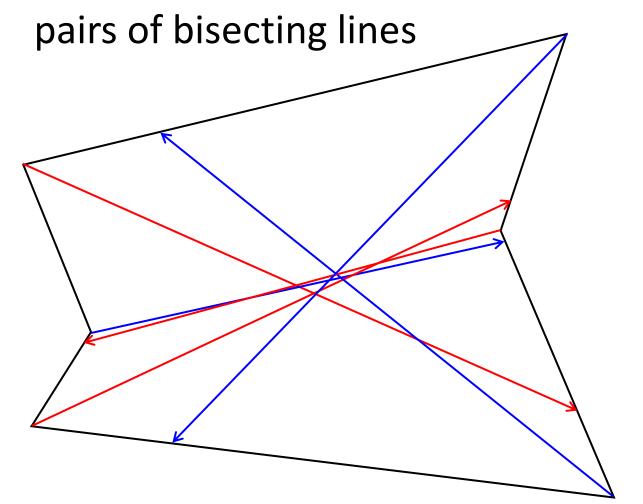
For triangle we had 
$$t = \left(1 + \sqrt{\frac{2-w}{1+w}}\right)^{-1}$$
;

For trapezoid we had 
$$t = \left(1 + \sqrt{\frac{2 - (1 - \gamma^2)w}{1 + \gamma^2 + (1 - \gamma^2)w}}\right)^{-1}$$
, for tilting parallel to base, top base = bottom base scaled by  $\gamma$ 

In general case things may not simplify so nicely (i.e. T(w) will involve u, x and y).

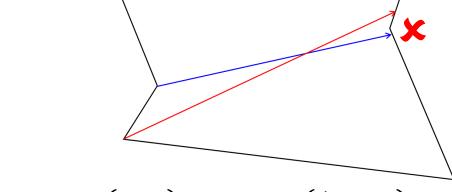


First step: Divide interior of polygon into sectors bordered by



Apply the bootlace formula trick at each vertex.

Pairs of lines defining a sector cannot 'include' a vertex because tilting won't work in this case.



In fact even u(x, y) = wx + (1 - w)y, won't work.

First step: Divide interior of polygon into sectors bordered by pairs of bisecting lines Q. Will this always cover an angle of  $\tau/2$  for bisection-convex polygons?