

Chapter 10. Section 3

The Recursive Euler Tour Algorithm

We proved that even degrees in a connected graph is sufficient to guarantee an Euler tour (the Euler-Hierholzer Theorem). Our inductive proof directly implies an algorithm to construct a tour.

Indeed, any inductive proof implies an algorithm: the size N case is constructed in the **induction step** from the size $N - 1$ case. This in turn is then constructed from the $(N - 1) - 1$ case. The reduction continues until we reach the **base case** which is constructed (proved) separately.

If we write the algorithm down explicitly it is naturally expressed using **recursion**: the algorithm is defined in terms of itself, but applied to a smaller input; in order to terminate, a smallest possible input (base case) is specified, for which the output is specified separately.

Here is a recursive algorithm for constructing an Euler tour:

Recursive_Euler_Tour(G)

Input: connected graph G in which all vertices have even degrees

if $E(G) = \emptyset$ **then return** $()$ **fi**; *# the base case: G is a single vertex; $()$ is the empty tour.*

$W :=$ a walk in G which forms a cycle, say, C ;

$G' := G \setminus C$, with connected components G_i , $k \geq 1$; *# $G \setminus C$ is G with the edges of C deleted.*

for $i := 1$ **to** k **do**

$S_i := \text{Recursive_Euler_Tour}(G_i)$; *# the recursive step.*

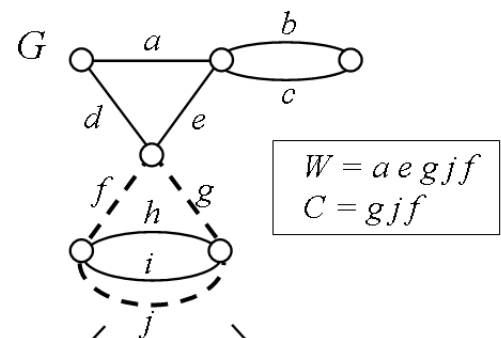
od;

$x_i :=$ a vertex in which G_i intersects C in G , $i = 1, \dots, k$;

$S := C$ with S_i (cycled to begin and end at x_i) inserted at x_i ;

return (S)

Example: in the graph G on the right the walk $W = a e g j f$ forms the cycle $C = g j f$



Deleting the edges of C gives a new graph ...

... with two connected components.

Each gives a new cycle.

Deleting these gives two new graphs ..

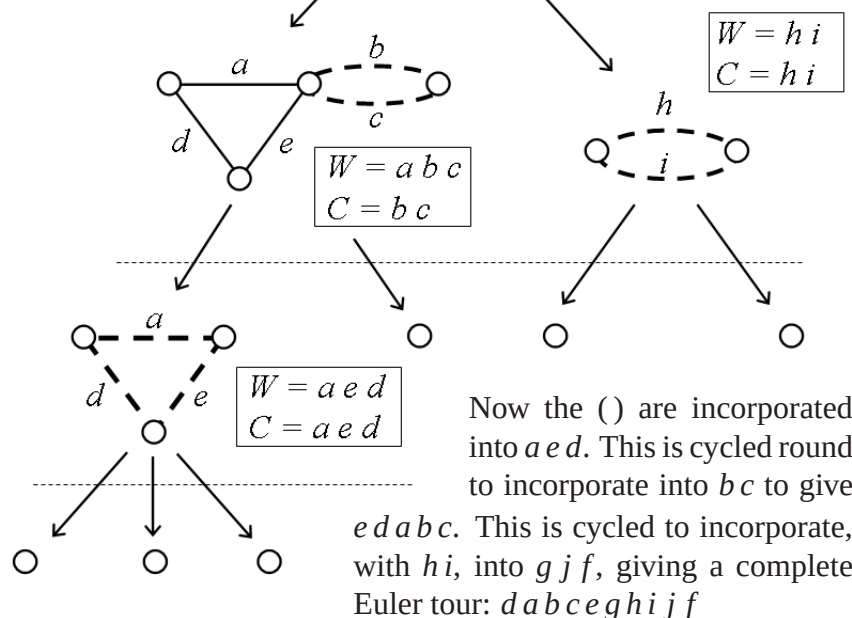
... each with two components

Three of this total of four components are trivial, with empty tours $()$.

The other has a new cycle.

Deleting this gives ...

... three trivial graphs, with empty tours $()$.



Now the $()$ are incorporated into $a e d$. This is cycled round to incorporate into $b c$ to give $e d a b c$. This is cycled to incorporate, with $h i$, into $g j f$, giving a complete Euler tour: $d a b c e g h i j f$