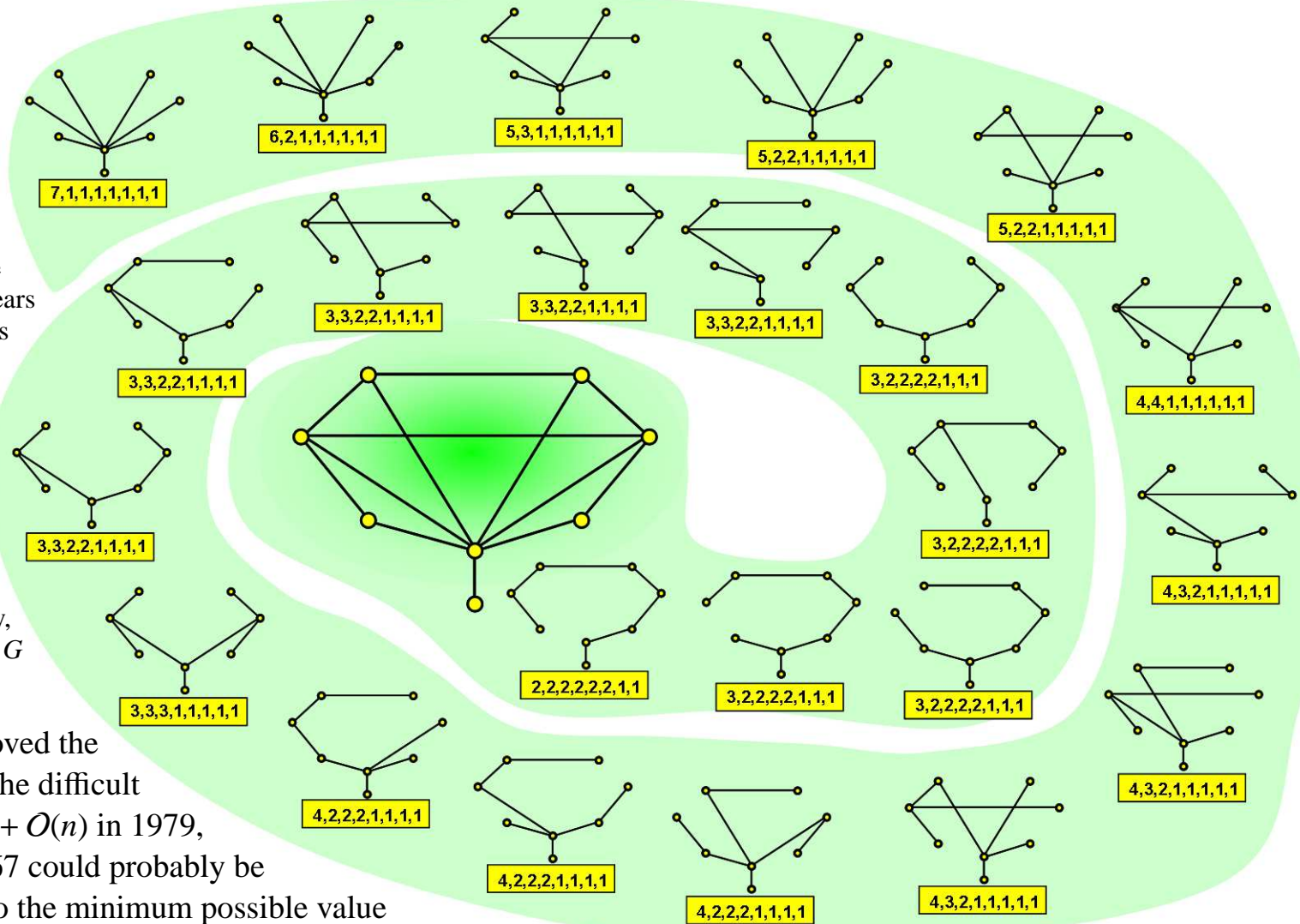


**The Panarboreal Formula** *Let  $\mathcal{T}_n$  denote the set of all unlabelled trees on  $n$  edges and denote by  $s(\mathcal{T}_n)$  the minimum number of edges which an  $(n+1)$ -vertex graph must have in order that it contains every tree in  $\mathcal{T}_n$  as a subgraph. Then  $s(\mathcal{T}_n) \sim cn \log n$  where  $c$  is a constant satisfying  $1/2 \leq c \leq 5/\log 4$ .*

There are 23 unlabelled trees having 7 edges; they are shown on the right, lexicographically ordered by degree sequence, together with an 8-vertex graph with 13 edges in which each one may be found as a subgraph. This appears to be the largest  $n$  for which  $s(\mathcal{T}_n)$  has been calculated exactly; but it is easy to establish that  $s(\mathcal{T}_n) \geq (1/2)n \log n$ . For, given  $k$ ,  $1 \leq k \leq n+1$  we may always choose a tree in whose degree sequence the  $k$ -th entry  $\geq n/k$ . But now the same must hold for the degree sequence of any graph  $G$  containing this tree. So if  $G$  contains each tree in  $\mathcal{T}_n$  and has degree sequence, say,  $(d_1, \dots, d_{n+1})$ , then number of edges in  $G = \frac{1}{2} \sum_{k=1}^{n+1} d_k \geq \frac{1}{2} \sum_{k=1}^{n+1} n/k > \frac{1}{2} n \log n$ .

Fan Chung and Ron Graham proved the easy lower bound on  $s(\mathcal{T}_n)$  and the difficult upper bound of  $(5/\log 4)n \log n + O(n)$  in 1979, mentioning that  $5/\log 4 \approx 3.6067$  could probably be improved, possibly even down to the minimum possible value of  $1/2$ . This challenge has yet to be met.



**Web link:** [math.ucsd.edu/~fan/](http://math.ucsd.edu/~fan/) (where all Chung's papers may be found in pdf format)

**Further reading:** *Erdős on Graphs: His Legacy of Unsolved Problems*, by Fan Chung and Ronald Graham, AK Peters, 1998, section 3.5.1.